

Dependent Types in the Idris Programming Language

Part 4: Type-level State Machines

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- Is it possible to parameterise a type by whether it is infinite or lazy?
- Could `covering` or `total` be default rather than `partial`?
- Where can we learn more about theorem proving in Idris?

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Questions from Lecture 3 (2)

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- What is the hardest/most annoying thing about implementing Idris?
- What's your favourite undocumented feature of Idris?

To create a server:

- Create a *socket*
- *Bind* the socket to a port
- *Listen* for connections to the socket
- *Accept* a connection:
 - This gives us a *new* socket for the connection
 - Continue *listening* on the original socket

Creating a server

```
int socket(int domain, int type, int protocol);

int bind(int socket, const struct sockaddr *address,
         socklen_t address_len);
int listen(int socket, int backlog);
int accept(int socket,
           struct sockaddr *restrict address,
           socklen_t *restrict address_len);

int connect(int socket, struct sockaddr *address,
            socklen_t address_len)
/* ... */
```

Which int is which? (Sockets)

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Which int is which? (Errors)

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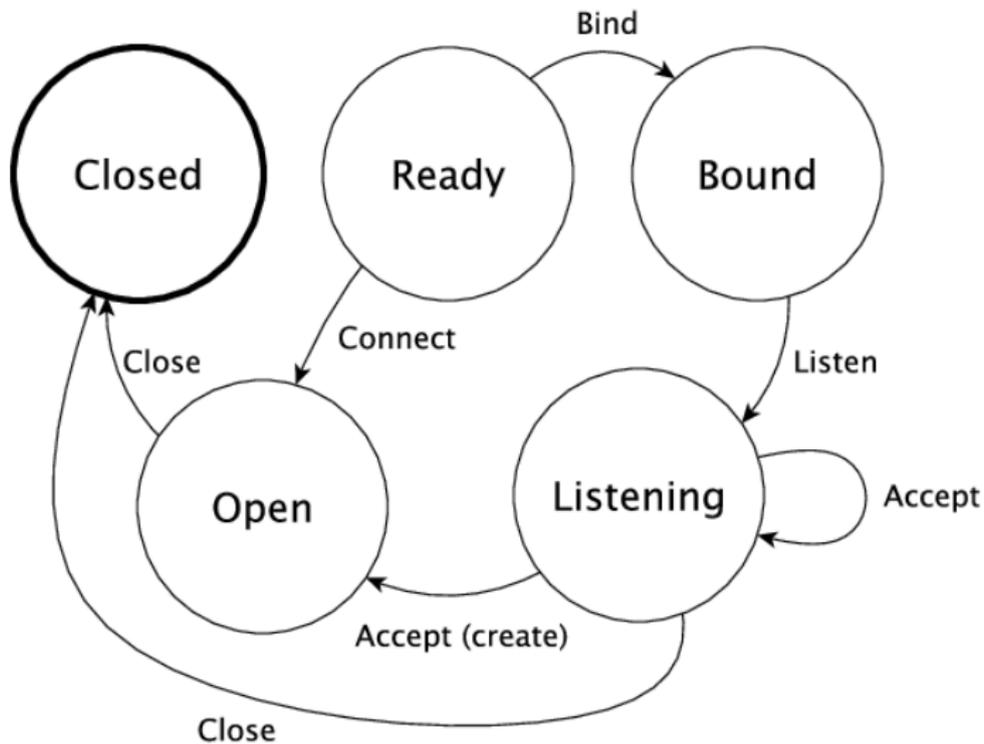
Creating a server

```
socket :: Family -> SocketType ->
        ProtocolNumber -> IO Socket

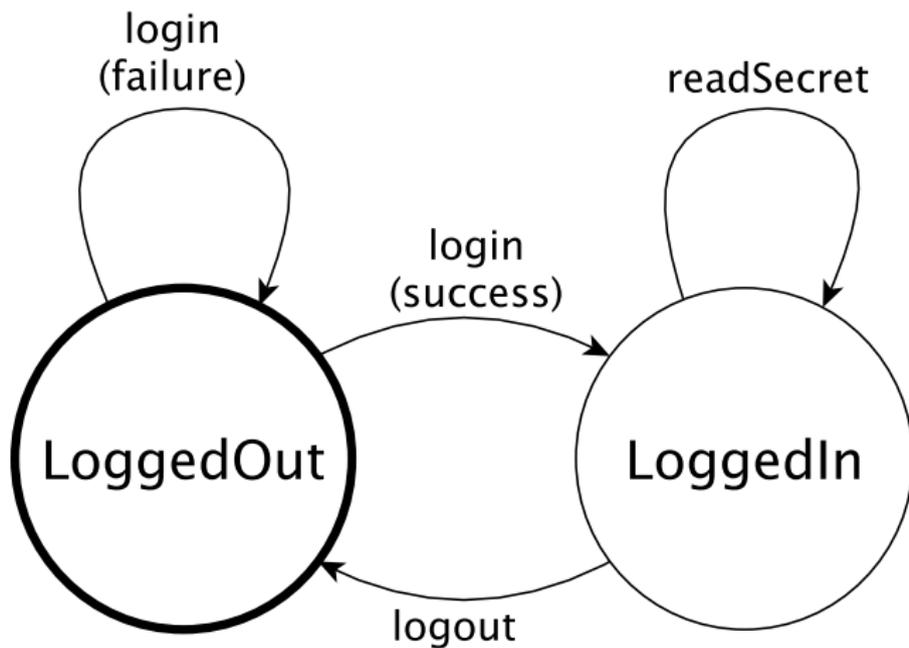
bind :: Socket -> SockAddr -> IO ()
listen :: Socket -> Int -> IO ()
accept :: Socket -> IO (Socket, SockAddr)

connect :: Socket -> SockAddr -> IO ()
```

Socket States



Smaller example: A Data Store



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Inspiration: *Separation logic, Ynot, Indexed Monads, Linear Types*

. . .

The type of stateful programs

```
ST : (m : Type -> Type) -- Underlying context
    -> (ty : Type)       -- Result type
    -> List (Action ty)  -- State transitions
                                -- may depend on the result
    -> Type
```

A DataStore Program

```
getData : (ConsoleIO m, DataStore m) => ST m () []
getData = do st <- connect
             OK <- login st
             | BadPassword => do putStrLn "Failure"
                                disconnect st

             secret <- readSecret st
             putStrLn ("Secret is: " ++ show secret)
             logout st
             disconnect st
```

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- *input resources*
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The type of stateful programs

```
STrans : (m : Type -> Type) -- Underlying context
        -> (ty : Type)       -- Result type
        -> Resources        -- Input states
        -> (ty -> Resources) -- Output states
                                -- may depend on the result
        -> Type
```

A resource

```
data Resource : Type where
  (:::) : Var -> Type -> Resource
```

Some example resource lists

```
[]
[door ::: Door Closed, count ::: State Int]
[d1 ::: Door Closed, d2 ::: Door Open,
 st ::: Store LoggedIn]
```

```
data Action : Type -> Type where
  Stable : Var -> Type -> Action ty
  Trans  : Var -> Type -> (ty -> Type) -> Action ty
  Remove : Var -> Type -> Action ty
  Add    : (ty -> Resources) -> Action ty
```

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```

Example actions

```
Stable st (Store LoggedIn)
Trans st (Store LoggedIn) (const (Store LoggedOut))
Remove st (Store LoggedOut)
Add (\var => [var :: Store LoggedOut])
```

```
data Action : Type -> Type where
  Stable : Var -> Type -> Action ty
  Trans  : Var -> Type -> (ty -> Type) -> Action ty
  Remove : Var -> Type -> Action ty
  Add    : (ty -> Resources) -> Action ty
```

Example actions, alternative notation

```
st ::: Store LoggedIn
st ::: Store LoggedIn :-> Store LoggedOut
remove st (Store LoggedOut)
add (Store LoggedOut)
```

General idea: Define *interfaces* for operations in **ST**:

- Abstract away underlying *resource types* (**Store** here)
- Implement interface for different *contexts*
- Use only needed interfaces

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The DataStore Interface

```
interface DataStore (m : Type -> Type) where
  ...

implementation DataStore IO where
  ....
```

Reading the secret

```
readSecret : (store : Var) ->  
            ST m String [store ::: Store LoggedIn]
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Logging out

```
logout : (store : Var) ->  
        ST m () [store ::: Store LoggedIn :->  
                Store LoggedOut]
```

Logging in (almost...)

```
login : (store : Var) ->  
  ST m () [store ::: Store LoggedOut :->  
           Store LoggedIn]
```

Logging in (with possible failure)

```
data LoginResult = OK | BadPassword

login : (store : Var) ->
  ST m Result
  [store ::: Store LoggedOut :->
    (\res => Store (case res of
      OK => LoggedIn
      BadPassword => LoggedOut))]
```

Connecting to a store

```
connect : ST m Var [add (Store LoggedOut)]
```

The DataStore Interface

```
interface DataStore (m : Type -> Type) where
  Store : Access -> Type
  connect : ST m Var [add (Store LoggedOut)]
  disconnect : (store : Var) ->
    ST m () [remove store (Store LoggedOut)]

{- ... login, logout, readSecret elided ... -}
```

To implement an interface, we need a *concrete* representation for `Store`:

The DataStore Interface

```
implementation DataStore IO where
  Store _ = State String

  {- ... -}
```

We also need the ability to *create*, *destroy*, *read* and *write* concrete resources.

We can create `new` resources, and `delete` resources, of type `State ty`:

Creating and destroying

```
new : (val : state) ->
      STrans m Var res
      (\lbl => (lbl ::: State state) ::: res)

delete : (lbl : Var) ->
         {auto prf : InState lbl (State st) res} ->
         STrans m () res (const (drop res prf))
```

We can **read** and **write** resources as long as we know they are of type **State ty**:

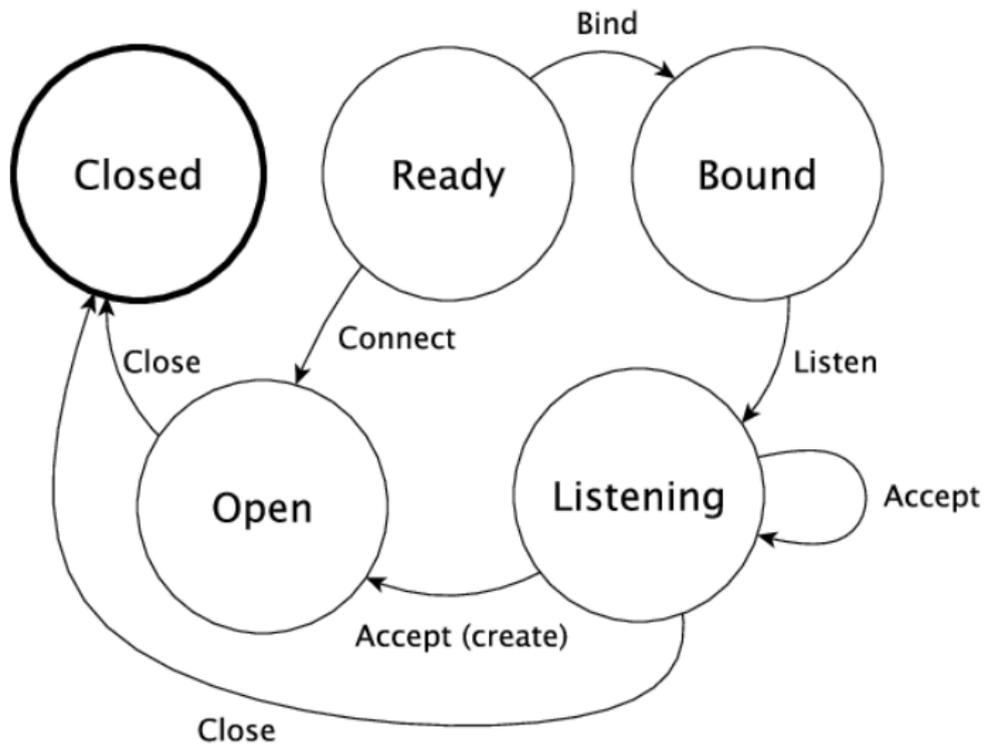
Reading and writing

```
read : (lbl : Var) ->
  {auto prf : InState lbl (State ty) res} ->
  STrans m ty res (const res)
write : (lbl : Var) ->
  {auto prf : InState lbl ty res} ->
  (val : ty') ->
  STrans m () res
    (const (updateRes res prf (State ty')))
```



Demonstrations: States in Action

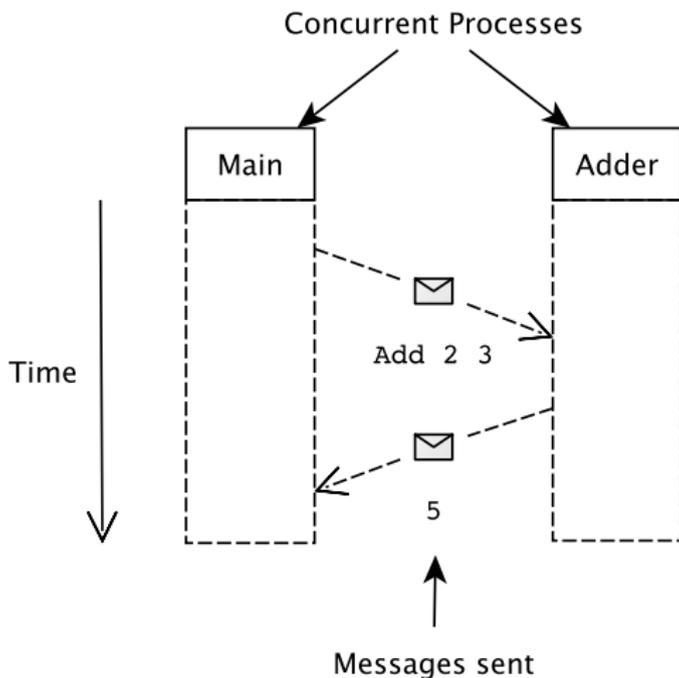
Socket States, Revisited



The Idris run-time system supports *message passing* concurrency

- A process can **spawn** another process
- A process can create a **Channel**, using:
 - **connect**, which initiates a connection to another process
 - **listen**, which waits for incoming connections
- Processes can **send** and **receive** messages on a **Channel**

Message Passing Concurrency in Idris



To write correct concurrent programs in this style, we'd like to ensure, at least:

- *Requests* (like `Add 2 3`) and *Responses* (like `5`) are well-typed w.r.t. each other
- *Server* processes (like `Adder`) run indefinitely
- Server processes always complete responses to requests

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A job for *Session Types*?

Session types:

- Describe the *state* of a communication channel
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ST can capture this! We can:

- Create threads
- Create channels for communicating between threads
- Give types to channels for tracking communication state

Channel States

```
data Actions : Type where
  Send : (a : Type) -> (a -> Actions) -> Actions
  Recv  : (a : Type) -> (a -> Actions) -> Actions
  Done  : Actions
```

Sending and Receiving

```
send : (chan : Var) -> (val : ty) ->  
      ST m () [chan ::: Channel (Send ty f) :->  
              Channel (f val)]  
  
recv : (chan : Var) ->  
      ST m ty [chan ::: Channel (Recv ty f) :->  
              (\res => Channel (f res))]
```



Demonstration: Concurrency with **ST**